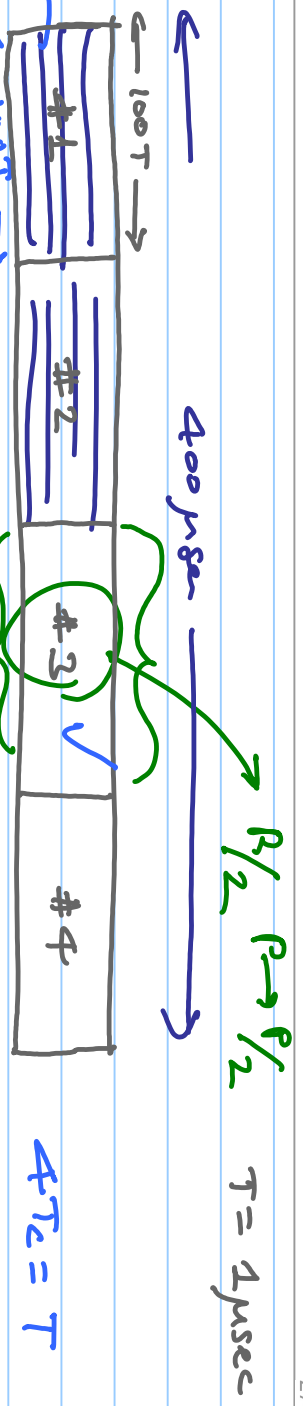


Recall:



TDMA

→ (max) per user bit rate : $\frac{100}{400\mu s} = 250 \text{ kbps} \leftarrow R$

→ 800 rate $\rightarrow 4 \times R$ $\underline{1000 \text{ kbps}}$ (1 Mbps)

DS-SSMA

→ Every user, in every loop will send 25 bits $\frac{25}{100\mu s} = 250 \text{ kbps}$

Sum rate = $4 \times R$ ✓

"No synch Banding" \rightarrow TDMA \rightarrow Repeat the bit

In DS-SSMA $N_f \leftarrow P$

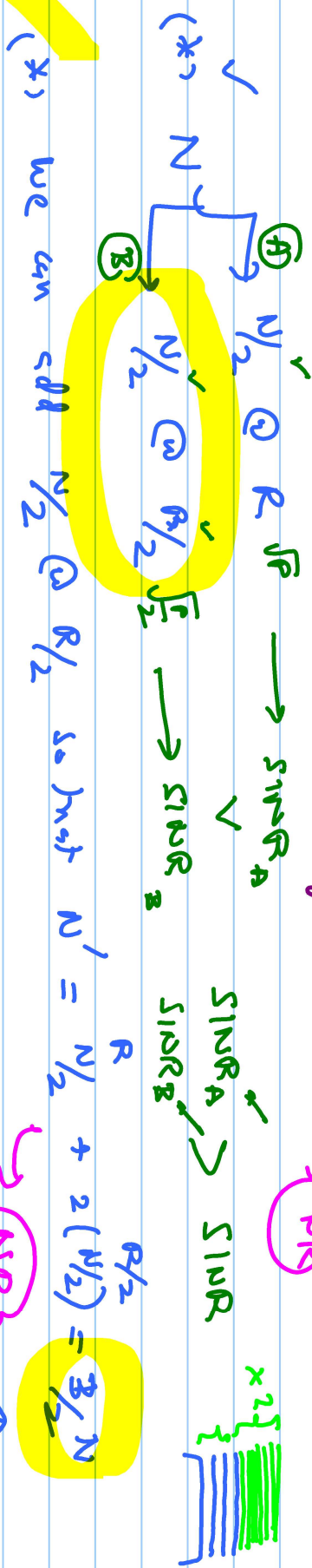
can all users @ rate R, power P

$$SINR = \frac{W \cdot P}{(W-1)P + \sigma^2}$$

target SINR

channel noise

$$W = \frac{T}{T_c}$$



$N/2 @ R; N/2 @ R/2; N/4 @ R$

$\rightarrow NR' \rightarrow \frac{SNR}{4} \leftarrow 25\%$

$P \rightarrow SINR \xrightarrow{A} SINR \xrightarrow{B} P'$

$P' > P$

Uplink Multiple Access

1. TDMA (unsprad) → single carrier comm. ✓

→ orthogonal support → single Rx in LOS links

✓ 'narrow banding' possible by repetition coding → no UL power control required

✗ however, the sum rate ↓

✗ optimal Rx for non-LOS channels (ISI channels → multiplex)
is MLSE ⇒ complexity scales exponentially with bandwidth
• M^{L-1} $M \rightarrow$ mod. size ; $L \rightarrow$ ISI $L \uparrow$

2. DS-SSMA

→ Pseudo orthogonal codes for LOS links ✗ → tight power control required on UL

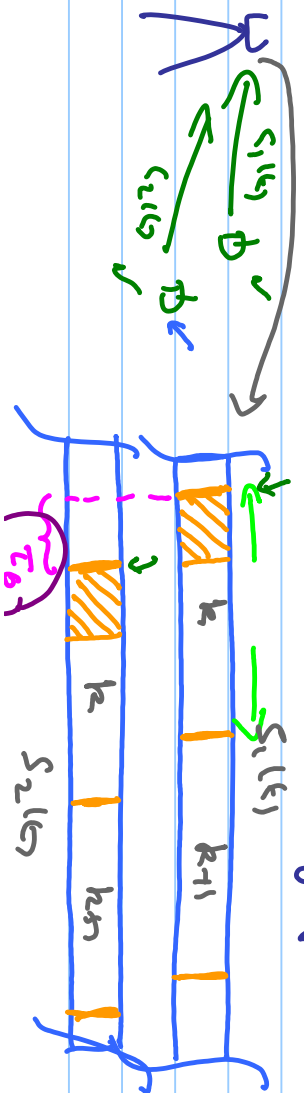
✓ Data ↓ with narrow-banding possible with statistical multiplexing
(i.e., sum rate does not ↓)

✓ Optimal Rx for ISI channel with interference user interference
is very complex MUD $M \uparrow$ $L \uparrow$ $L \rightarrow$ ISI $L \uparrow$

3. OFDM / OFDMA

- orthogonal subcarriers
- strict power control not reqd.
- UL resource allocation
- optimal Rx for ISI channels is simple →
- complexity scales only linearly with BW

Question: How to prevent orthogonality even when the codes have varying errors?



$$s_1(t) = \sum_k c_k(t)$$

$$s_2(t) = \sum_k c_k(t)$$

$$r(t) = s_1(t) + s_2(t - t_0)$$

→ Extend the length of the codes by say T_{cp} CP → cyclic prefix
 so that $T = T_u + T_{cp}$ "sumtime"

→ As long as $0 \leq \tau_0 < T_{cp}$, we can show that $s_1(t) & s_2(t - \tau_0)$ can be separated orthogonally. "useful time"

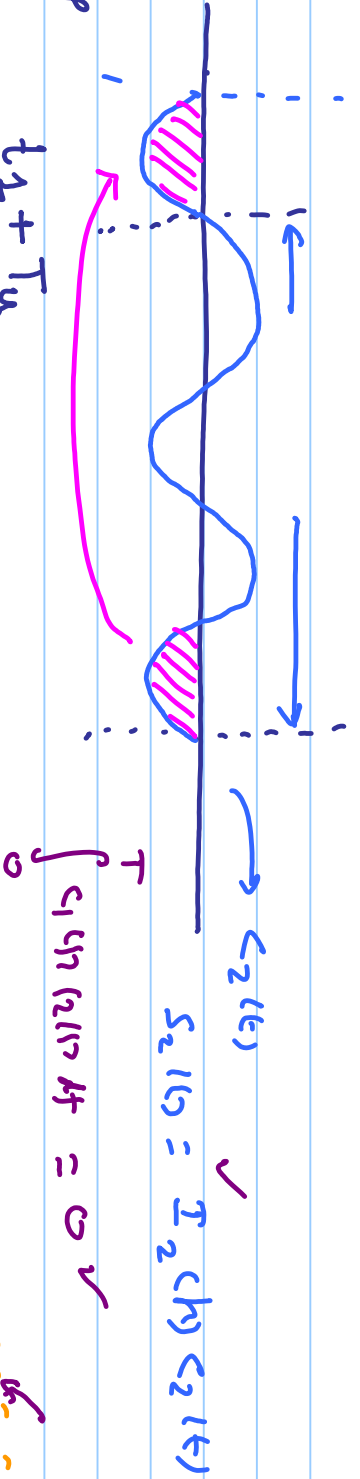
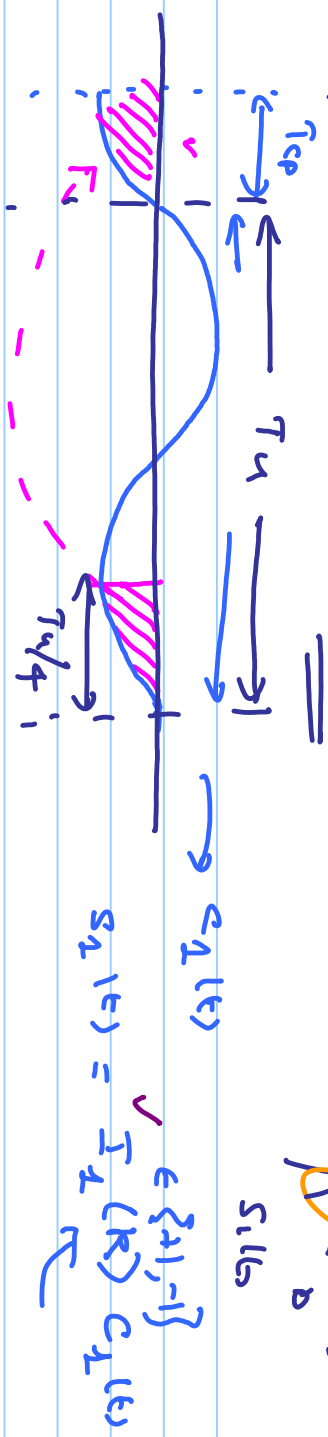
WTC $\sum_{i=1}^N$

Key Idea: Check orthogonality only over T_u (out of T) seconds. (in T_{cp})

DFDM/OFDM \Rightarrow in discrete time, check orthogonality for $\sum_{n=0}^{N-1} s(n)$ only N samples out of $N + N_{cp}$ received samples $\sum_{n=0}^N$

$s_2(t - \tau_0)$

eg: \checkmark Reed sinusoids \rightarrow codes



$$\int_{t_1}^{t_1 + \tau_u} c_1(t) c_2(t - \tau_0) dt = 0, \quad \checkmark \text{ provided } t_1 > \tau_0 \quad \checkmark$$

$$\int_0^T c_1(t) c_2(t) dt = 0, \quad \checkmark \quad \text{no } |B|$$

\checkmark $0 \leq \tau_0 < \tau_{cp}$ \checkmark

Recall: Trigonometric Identities $f_0 = \frac{1}{T}$;

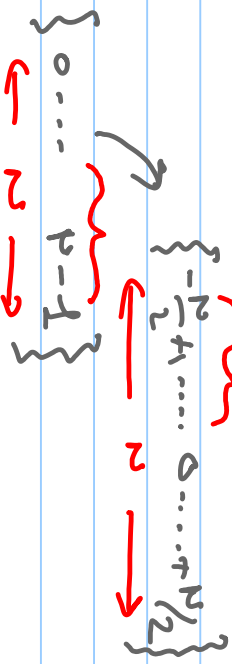
$$(*) \int_0^T \sin 2\pi f_0 t \sin 4\pi f_0 t \, dt \rightarrow \frac{1}{2} \left[\cancel{\cos(\pi f_0 T)} + \cancel{\cos(\pi f_0 T)} \right] = 0$$

$$(*) \int_0^T \sin 2\pi f_0 t \sin(4\pi f_0 t - \underbrace{\pi f_0 T}_{=0}) \, dt = 0 ;$$

codes

From real sinusoids

(harmonic codes not labeled)



(*) Discrete-time ✓

(*) Complex symbols ✓

(*) "Block of samples" ✓

$$-j \left[e^{j \frac{2\pi}{N} nm} \right], n, m \in \{0, \dots, N-1\}$$

DFT basis set

DFT

DFT

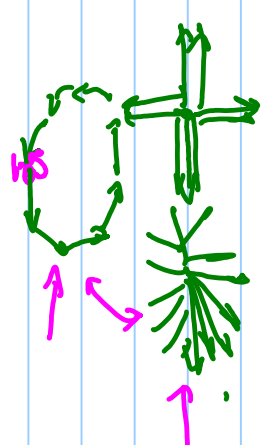
"code-set"

$$q_n = \begin{bmatrix} e^{j\frac{2\pi}{N}n \cdot 0} \\ e^{j\frac{2\pi}{N}n \cdot 1} \\ \vdots \\ e^{j\frac{2\pi}{N}n \cdot (N-1)} \end{bmatrix} \xrightarrow{\text{IDFT}} F = \frac{1}{\sqrt{N}} \begin{bmatrix} 1 & 1 & \dots & 1 \\ q_0 & q_1 & \dots & q_{N-1} \\ \vdots & \vdots & \ddots & \vdots \\ q_{N-1} & q_1 & \dots & q_{N-1} \end{bmatrix} \xrightarrow{\text{DFT}} F^H = \frac{1}{\sqrt{N}} \begin{bmatrix} -q_0^* & -q_1^* & \dots & -q_{N-1}^* \\ q_0 & q_1 & \dots & q_{N-1} \\ \vdots & \vdots & \ddots & \vdots \\ q_{N-1} & q_1 & \dots & q_{N-1} \end{bmatrix}$$

$k \rightarrow$ Block #
 $n \rightarrow$ Sub-carrier (code) index
 $m \rightarrow$ Time sample.

Prep #0 :

$$\frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} e^{j\frac{2\pi}{N}nm} = \begin{cases} 0, & \text{otherwise} \\ \sqrt{N}, & m=0 \end{cases} \quad (n=1 \dots N-1)$$



Prep #1 :

$$\langle \bar{q}_m, \bar{q}_n \rangle = \bar{q}_n^H \bar{q}_m = \frac{1}{\sqrt{N}} \sum_{m=0}^{N-1} e^{-j\frac{2\pi}{N}nm} e^{j\frac{2\pi}{N}ms} = \sqrt{N} \delta_{nm}$$

Delay $l=0$

Prop # 2 :

Delay $L=0$

$$\begin{aligned} &\langle \bar{r}_n, \bar{r}_{n'} \rangle \quad n' \neq n, \quad n, n' \in \{0, 1, \dots, N-1\} \\ &= \frac{1}{\sqrt{N}} \sum_{m=0}^{N-1} e^{j\frac{2\pi}{N} (n'-n)m} \\ &= 0; \quad n'-n \in \{ \pm 1, \pm 2, \dots, \pm (N-1) \} \end{aligned}$$

Prop # 3 :

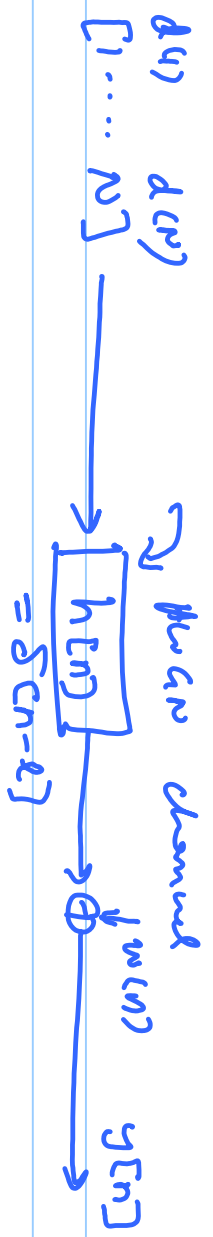
Delay $L \leq N/2$

$$\begin{aligned} &\langle \bar{r}_n, \bar{r}_{n(L)} \rangle \\ &= \frac{1}{\sqrt{N}} \sum_{m=0}^{N-1} e^{-j\frac{2\pi}{N} nm} e^{j\frac{2\pi}{N} n (m-L)} \\ &= \frac{1}{\sqrt{N}} \cdot e^{-j\frac{2\pi}{N} nL} \sum_{m=0}^{N-1} e^{j\frac{2\pi}{N} n m} \end{aligned}$$

for each "n", we get a different exp.

Prop # 4 :

$$\begin{aligned} &\langle \bar{r}_n, \bar{r}_{n(L)} \rangle \\ &= \frac{1}{\sqrt{N}} e^{-j\frac{2\pi}{N} nL} \sum_{m=0}^{N-1} e^{j\frac{2\pi}{N} n (n'-n)} \\ &= 0; \end{aligned}$$

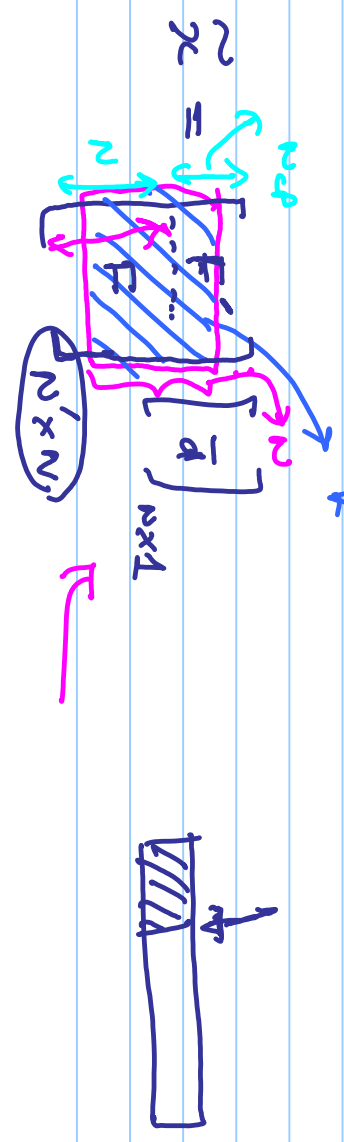


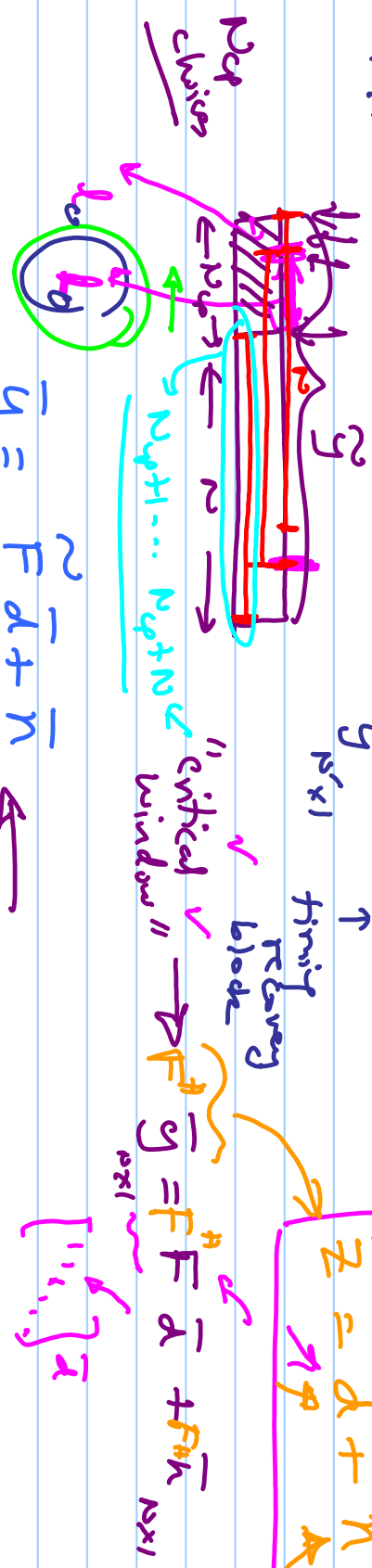
$$x(n) \quad \bar{x} \quad \begin{matrix} N \times 1 \\ N \times 1 \end{matrix} = \frac{1}{\sqrt{N}} \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ \bar{g}_0 & \bar{g}_1 & \bar{g}_2 & \dots & \bar{g}_{N-1} \\ 1 & 1 & 1 & \dots & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 1 \\ \dots \\ 1 \end{bmatrix} \begin{bmatrix} d(k_1) \\ d(k_2) \\ \vdots \\ d(k_N) \end{bmatrix}$$

IDFT \rightarrow

$$\bar{x} = F d$$

(*) Adding CP $\bar{x} \rightarrow x$
 $N+1 \times 1$
 $N \times 1$



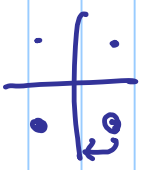


"no channel estimation is required"

$$\tilde{y} = F^H F \tilde{x} + F^H \tilde{n}$$

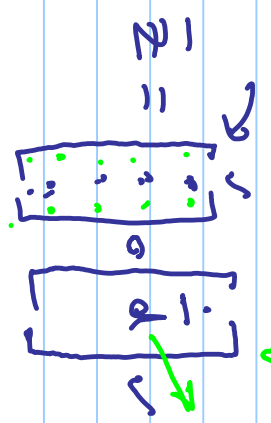
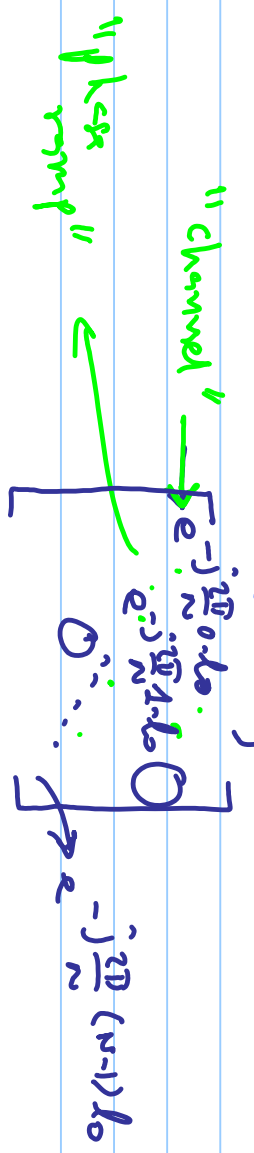
$$\begin{bmatrix} 1 \\ \vdots \\ \vdots \end{bmatrix} \tilde{x}$$

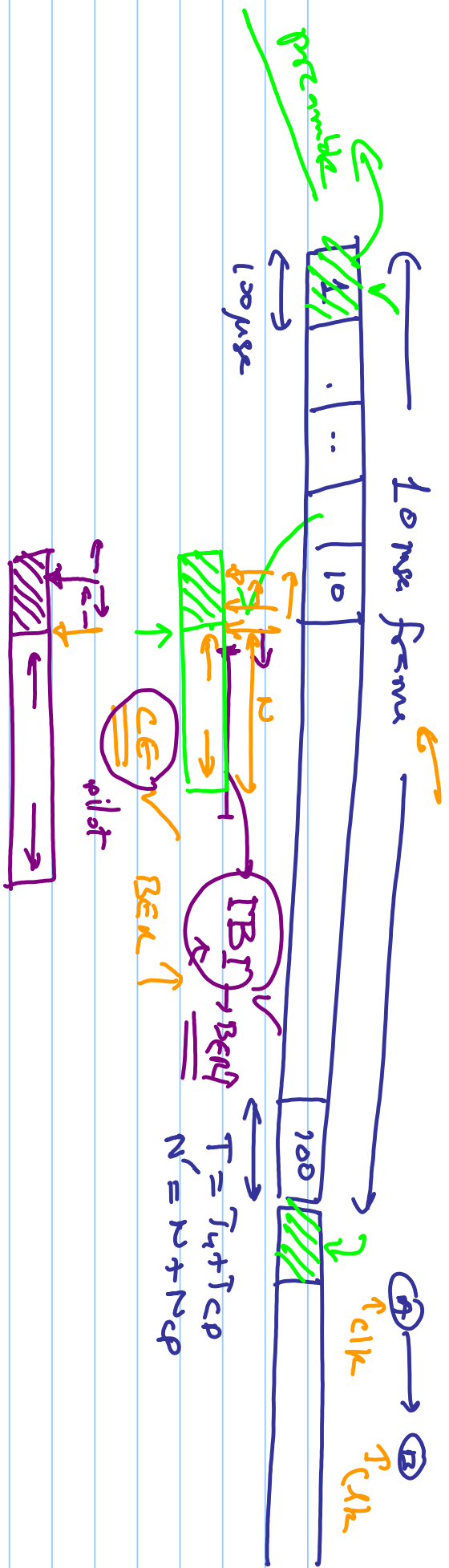
$$F^H F = I$$



$$\tilde{y} = F \tilde{x} + \tilde{n}$$

$$F^H \tilde{y} = F^H F \tilde{x} + F^H \tilde{n}$$





$$N = N_c * h \rightarrow ?$$